A tale of three disciplines

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Let N(t) be te worst possible number of trajectories compatible with an observation of length t A network is trackable if N(t) grows subexponentially

[Crespi et al. 05]

 $N(t) \approx 0$

Here: number of possibilities asymptotically zero

➔ Trackable



Worst case : RRRRR... →

 $N(t) \approx t$

Polynomial number of possibilities





Worst case : RGRGRG...→

 $N(t) \approx 2^{t/2}$

Exponential number of possibilities

➔ Not trackable

Trackability : the formal problem



For each possible color, we define the corresponding matrix by erasing the incompatible columns from A:

$$A_r = \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} \quad A_g = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{pmatrix}$$

Trackability : the formal problem

To a given observation, associate the corresponding product:

$$A_{r} = \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} A_{g} = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{pmatrix}$$
$$A_{r}A_{g}A_{r} = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix}$$

The number of possible trajectories is given by the sum of the entries of the matrix

Outline

• Matrix semigroups and some tools to analyze them

• LMI methods for switching systems stability

• Applications

• Conclusion and perspectives

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Switching systems

$$\mathbf{x}_{t+1} = \begin{array}{c} \mathbf{A}_0 \ \mathbf{x}_t \\ \mathbf{A}_1 \ \mathbf{x}_t \end{array}$$

Point-to-point Given x_0 and x_* , is there a product (say, $A_0 A_0 A_1 A_0 \dots A_1$) for which $x_*=A_0 A_0 A_1 A_0 \dots A_1 x_0$?

Mortality Is there a product that gives the zero matrix?

Boundedness Is the set of all products $\{A_0, A_1, A_0A_0, A_0A_1, ...\}$ bounded?

$\mathbf{x}_{t+1} = \begin{array}{l} \mathbf{A}_{0} \mathbf{x}_{t} \\ \mathbf{A}_{1} \mathbf{x}_{t} \end{array}$

Global convergence to the origin Do all products of the type $A_0 A_0 A_1 A_0 \dots A_1$ converge to zero?

The spectral radius of a matrix A controls the growth or decay of powers of A

$$\rho(A) = \lim_{t \to \infty} ||A^t||^{1/t}$$
 The powers of A converge to zero iff $\rho(A) < 1$

The joint spectral radius of $\Sigma = \{A_0, A_1\}$ is given by $\rho(\Sigma) = \lim_{t \to \infty} \max_{A_i \in \Sigma} ||A_1 A_2 \dots A_t||^{1/t}$

All products of A₀ and A₁ converge to zero iff $\rho(\Sigma) < 1$















$$\rho(\Sigma) = \lim_{t \to \infty} \left[\max_{A_i \in \Sigma} ||A_1 A_2 \dots A_t|| \right]^{1/t}$$

The joint spectral radius answers the **stability** problem

$$\check{\rho}(\Sigma) = \lim_{t \to \infty} \left[\min_{A_i \in \Sigma} ||A_1 A_2 \dots A_t|| \right]^{1/t}$$

The joint spectral subradius answers the **stabilizability** problem

$$\rho_{p}(\Sigma) = \lim_{t \to \infty} \left[m^{-t} \sum_{A_{i} \in \Sigma^{t}} \|A_{1}A_{2} \dots A_{t}\|^{p} \right]^{1/(pt)}$$
The p-radius addresses
the... p-weak stability
[J. Protasov 10]

$$\bar{\rho}(\Sigma) = \lim_{t \to \infty} \left[\prod_{A_{i} \in \Sigma^{t}} \|A_{1}A_{2} \dots A_{t}\| \right]^{1/(tm^{t})}$$
The Lyapunov exponent
answers the stability
with probability one
(Cfr. Oseledets Theorem)

The joint spectral characteristics: Mission Impossible?

Theorem Computing or approximating ρ is NP-hard

Theorem The problem $\rho \leq 1$ is algorithmically undecidable

Conjecture The problem ρ <1 is algorithmically undecidable





See

Theorem Even the question « $|\check{
ho} - r| \le a + b\check{
ho}$?» is algorithmically undecidable for all (nontrivial) a and b

Theorem The same is true for the Lyapunov exponent

Theorem The p-radius is NP-hard to approximate

[Blondel Tsitsiklis 97, Blondel Tsitsiklis 00, J. Protasov 09]

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LMI methods

$$\rho(\Sigma) = \lim_{t \to \infty} \left[\max_{A_i \in \Sigma} ||A_1 A_2 \dots A_t|| \right]^{1/t}$$

• Idea: look for a vector norm such that $\forall A \in \Sigma, \forall x : |x| \le 1, |Ax| \le r$



Boundedness (and $ho \leq r$)

SDP methods

• Theorem For all $\epsilon > 0$ there exists a norm such that

 $\forall A \in \Sigma, \forall x, |Ax| \leq (\rho + \epsilon) |x| \qquad \text{[Rota Strang, 60]}$

• John's ellipsoid Theorem: Let K be a compact convex set with nonempty interior symmetric about the origin. Then there is an ellipsoid E such that $E \subset K \subset \sqrt{nE}$



SDP methods

• Theorem The best ellipsoidal norm $\|\cdot\|_{E_*}$ approximates the joint spectral radius up to a factor \sqrt{n} [Ando Shih 98]



• Theorem The best ellipsoidal norm of a set of m (nonnegative) matrices also approximates the joint spectral radius up to a factor \sqrt{m}

 $ho \leq \max ||A||_{E_*} \leq \sqrt{m}
ho$ [Blondel Nesterov 05]

LMI methods

• **Problem** How to compute the best ellipsoidal norm?

$$\begin{aligned} &\inf_{r \in \mathbb{R}^+} & r \\ &\text{s.t.} \\ &A^T P A & \preceq & r^2 P, \quad \forall A \in \Sigma \\ &P & \succeq & 0. \end{aligned}$$

$$\Leftrightarrow \frac{|Ax|_P}{|x|_P} \leq r$$

- This is « just » a sufficient condition for stability
- Computable in polynomial time (interior point methods)
- The « CQLF method » (see [Mason Shorten 04]...)

LMI methods

- Other methods using LMIs have been proposed
 - Based on symmetric algebras [Blondel Nesterov 05]
 - Based on sum of squares [Parrilo Jadbabaie 08]

 All these methods can be seen as trying to approximate the best norm with more and more complex curves

...Or, the symmetric algebras (the SOS) allow for lifting the matrices in a higher-dimensional (but more gentle) space

There exist **algorithms** that approximate the joint spectral radius of arbitrary sets of n by n matrices up to an arbitrary accuracy ϵ in $O(n^{k\frac{1}{\epsilon}})$ operations

Yet another LMI method

• A strange semidefinite program

$$\min_{r \in \mathbb{R}^+} \qquad r$$
s.t.

$$\begin{array}{ccc} A_1^T P_1 A_1 & \preceq & r^2 P_1, \\ A_2^T P_1 A_2 & \preceq & r^2 P_2, \\ A_1^T P_2 A_1 & \preceq & r^2 P_1, \\ A_2^T P_2 A_2 & \preceq & r^2 P_2, \\ P & \succeq & 0. \end{array}$$

 $\rho \leq r$

[Goebel, Hu, Teel 06]

But also... [Daafouz Bernussou 01]
 [Bliman Ferrari-Trecate 03]
 [Lee and Dullerud 06] ...

Yet another LMI method

• An even stranger program:

 $\min_{r \in \mathbb{R}^+} \qquad r$ s.t. $A_1^T P A_1 \qquad \preceq \quad r^2 P,$ $(A_2 A_1)^T P (A_2 A_1) \qquad \preceq \quad r^4 P,$ $(A_2^2)^T P (A_2^2) \qquad \preceq \quad r^4 P,$ $P \qquad \succeq \quad 0.$



[Ahmadi, J., Parrilo, Roozbehani10]

Yet another LMI method

- Questions:
 - Can we characterize all the LMIs that work, in a unified framework?
 - Which LMIs are better than others?
 - How to prove that an LMI works?
 - Can we provide converse Lyapunov theorems for more methods?

From an LMI to an automaton

• Automata representation Given a set of LMIs, construct an automaton like this: A_1



- Definition A labeled graph (with label set A) is path-complete if for any word on the alphabet A, there exists a path in the graph that generates the corresponding word.
- Theorem If G is path-complete, the corresponding semidefinite program is a sufficient condition for stability. [Ahmadi J. Parrilo Roozbehani 11]

An obvious question: are there other Theorem valid criteria?



If G is path-complete, the corresponding semidefinite program is a sufficient condition for stability.

- Are all valid sets of equations coming from path-complete graphs?
- ...or are there even more valid LMI criteria?

Are there other valid criteria?

• Theorem Every non path-complete set of equations is not a sufficient condition for stability. [J. Ahmadi Parrilo Roozbehani 12]



• Corollary:

It is PSPACE complete to recognize sets of equations that are a sufficient condition for stability

These results are not limited to LMIs, but apply to other families of conic inequalities

What about the other quantities?

| | Arbitrary approximation | Arbitrary approximation in polynomial time | Arbitrary approximation for positive matrices | Decidability | |
|--------------------------------|----------------------------|--|---|--------------|--|
| Joint Spectral Radius | * | * | * | ? | |
| Joint Spectral Subradius | × | K | * | K | |
| Lyapunov Exponent | 8 | × | v | 8 | |
| p-radius | Depends on p | Depends on p | v | ? | |

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multihop control networks



How to model failures? LTIs with switched delays



$$x(t+1) = Ax + Bu(t-d)$$

How to model failures? LTIs with switched delays



$$x(t+1) = Ax + Bu(t-d_2)$$

$$A = \begin{pmatrix} 0 & 2 \\ 2 & 0 \end{pmatrix}, \quad B = \begin{pmatrix} 0 & 1 \end{pmatrix}^T$$
$$D = \{0, 1\}, \quad \sigma(t) = t \mod 2$$

LTIs with switched delays The linear controller



$$\begin{aligned} x(t+1) &= Ax(t) + Bu(v(t-d_{max}:t), \sigma(t-d_{max}:t)) \\ u_d(t) &= \sum_{t' < t: t' + \sigma(t') = t+d} v(t') \end{aligned}$$

LTIs with switched delays The linear controller





LTIs with switched delays





• Corollary

For both models there is a **PTAS** for the stability question:

for any required accuracy, there is a polynomial-time algorithm for checking stability up to this accuracy

Previous sufficient conditions for stability in [Hetel Daafouz Iung 07, Zhang Shi Basin 08]

• However:

Theorem the very stability problem is NP-hard

• Open questions

What about the stabilizability? What about the stability with probability one? (sufficient conditions in, e.g.

[Weiss et 09])

[[]J. D'Innocenzo Di Benedetto 12]

Design of LTIs with switched delays Delay dependent controller:

 Theorem for n=m=1, there is an explicit formula for a controller that achieves finite time stabilization.

(based on a generalization of the Ackermann formula for delayed LTI)

Example

$$A = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad B = \begin{pmatrix} 0 & 1 \end{pmatrix}^{T}$$

$$D = \{0, 1\}, \quad \sigma(t) = t \mod 2$$

$$x_{0} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$$x_{1} = A \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$x_{2} = A^{2} \begin{pmatrix} 1 \\ 0 \end{pmatrix} + Bv(1) + Bv(2) = \begin{pmatrix} 1 \\ 0 \end{pmatrix} + \begin{pmatrix} 0 \\ v(1) + v(2) \end{pmatrix}$$

uncontrollable

Design of LTIs with switched delays the delay-independent case

• Example: n=m=1, $D=\{0,1\}$: find k_1 , k_2 such that

$$\begin{aligned} x(t+1) &= ax(t) + (1 - \sigma(t)) & (k_1 x(t) + k_2 x(t-1)) \\ &+ \sigma(t-1) & (k_1 x(t-1) + k_2 x(t-2)) \end{aligned}$$

is stable, whatever switching signal $\sigma(t)$ occurs?

- a>3: impossible
- a=2: possible, but one needs memory
- a=1.1: possible with $k_2 = 0$
- a<1: no need for a controller!
- → Question: how to algorithmically decide if stabilization is possible, even in this simplest case?

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To a given observation, associate the corresponding product:

$$A_{r} = \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} A_{g} = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{pmatrix}$$
$$A_{r}A_{g}A_{r} = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix}$$

The number of possible trajectories is given by the sum of the entries of the matrix



The maximal total number of possibilities is

$$N(t) = \max\left\{ \left\| A \right\|_1 : A \in \Sigma^t \right\}$$

We are interested in the asymptotic worst case :

$$\lim_{t \to \infty} N(t)^{1/t} = \lim_{t \to \infty} \max\left\{ \left\| A \right\|_{1}^{1/t} : A \in \Sigma^{t} \right\}$$

This is a joint spectral radius!



The network is trackable iff

[Crespi et al. 05]

Theorem It is possible to check trackability in polynomial time

[J. Protasov Blondel 08]

 $\rho \leq 1$

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TCP Congestion control



Congestion avoidance strategy: the **AIMD algorithm**

(Additive Increase Multiplicative Decrease)

If the node i does not see congestion $w_i(t+1) = w_i(t) + \alpha$

If the node i sees congestion

$$w_i(t+1) = \beta w_i(t), \ \beta < 1$$

AI

MD

TCP Congestion control



Congestion avoidance strategy: the **AIMD algorithm**

(Additive Increase Multiplicative Decrease)

This is a switching system!

[Scholte, Berman, Shorten 12]

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 - Other applications

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[Hernandez-Varga colaneri Middleton Blanchini 10]

Applications



1001011001101001...

[Blondel Cassaigne J. 09] [Rigo Berthe 10]

> [Shorten et al. 07] [Goebel, Sanfelice, Teel 09] [J. 09]

 $\operatorname{cap}(D) = \lim_{n \to \infty} \frac{\log_2 \delta_n(D)}{n}$

 $\operatorname{cap}(D) = \log_2 \rho(\Sigma(D))$

10001110101

[Dima Asarin 10]

Siegel 12]

[Moision Orlitsky Siegel 01]

Beal crochemore Moision

[Blondel J. Protasov 06]

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Conclusion and perspectives

- Many open questions and Conjectures:
 - For computer scientists:

Is ` ρ <1' decidable? [Blondel Megretski 05]

Write protocols for optimal conrol of computer networks

- For mathematicians: The finiteness conjecture
- For control theorists: ^{[Lagarias, Wang 95], [Bousch Mairesse 02] [J. Blondel 08]} What are the best control schemes for LTIs with switched delays?

How to design switching systems? [Lee Dullerud 06, Hetel Daafouz Iung 07, Zhang Shi Basin 08] [J. D'Innocenzo Di Benedetto 12]

 Meta-conclusion: many hybrid systems boil down to the study of switching systems and matrix semigroups, for which nice mathematical tools exist

Thanks!

Questions?

Ads

<u>The JSR Toolbox:</u> <u>http://www.mathworks.com/matlabcentral/</u> <u>fileexchange/33202-the-jsr-toolbox</u>

References: http://perso.uclouvain.be/raphael.jungers/

Joint work with A.A. Ahmadi (IBM-Watson), M-D di Benedetto (DEWS), V. Blondel (UCLouvain), A. Cicone (l'Aquila), A. D'innocenzo (DEWS), N. Guglielmi (l'Aquila), P. Parrilo (MIT), V. Protasov (Moscow), M. Roozbehani (MIT)...

