Fast oriented bounding box optimization on the rotation group $SO(3,\mathbb{R})$

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Coming up next...

The problem: minimum-volume OBB

Exact methods in 2D and 3D Classical approaches for the 3D case Our goal

Bringing optimization into the game

How to solve an optimization problem?

Results

Conclusion

The problem

Given a set of n points \mathcal{X} in 3D, find the minimum-volume arbitrarily oriented bounding box enclosing \mathcal{X} .

Collision detection, intersection tests, object representation, data approximation... (BV trees...)

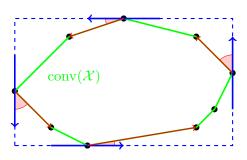
In 2D: the rotating calipers method

A minimum-area rectangle circumscribing a convex polygon has at least one side flush with an edge of the polygon.

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Compute the convex hull: $\mathcal{O}(n \log n)$



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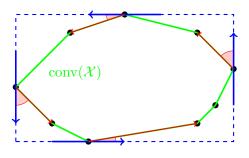
A minimum-area rectangle circumscribing a convex polygon has at least one side flush with an edge of the polygon.

Compute the convex hull:

$$\mathcal{O}(n \log n)$$

Loop on all edges:

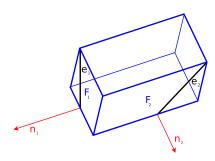
 $\mathcal{O}(n) \to \mathsf{easy}$ and efficient



A minimum-volume box circumscribing a convex polyhedron has at least one face flush with a face of the polyhedron?

In 3D: generalization of the rotating calipers?

A minimum-volume box circumscribing a convex polyhedron has at least one face two adjacent faces flush with a face edges of the polyhedron. [O'Rourke, 1985]

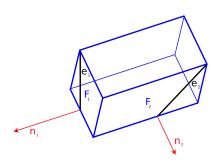


In 3D: generalization of the rotating calipers?

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Problem:

Loop on all pairs of edges and rotate the box while keeping edges flush $\rightarrow \mathcal{O}(n^3)$ time complexity...



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 - Quaranteed quality approximation methods: same problem...

What do we want?

Goal:

- Very good accuracy: find an optimal OBB in (nearly?) all cases
- ♦ If a suboptimal solution is returned, it should be close to the best one
- Computational cost has to be low

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Our approach: iterative algorithm based on optimization methods

Coming up next...

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Bringing optimization into the game

OBB fitting as an optimization problem Requirements
Going hybrid

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 $\begin{array}{c} \min \\ \text{over size, position, orientation} \\ \text{so that} \end{array} \quad \begin{array}{c} \text{the volume of the bounding box} \\ \text{all the points are in the box} \end{array}$

 $\begin{array}{c} \min \\ \text{over } \Delta \in \mathbb{R}^3, \text{ position, orientation} \\ \text{so that} \end{array} \qquad \begin{array}{c} \text{the volume of the bounding box} \\ \text{all the points are in the box} \end{array}$

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Smooth but constrained optimization problem

Unconstrained formulation

$$\min_{R \in SO(3,\mathbb{R})} \underbrace{\left(\begin{array}{cc} \min_{\Delta \in \mathbb{R}^3, \Xi \in \mathbb{R}^3} & \Delta_{\xi} \Delta_{\eta} \Delta_{\zeta} \\ \text{s.t.} & -\frac{1}{2} \Delta \leq R \mathbf{X}_i - \Xi \leq \frac{1}{2} \Delta \ \forall i = 1, \dots, N \end{array} \right)}_{f(R)}$$

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Unconstrained but non-differentiable optimization problem

... a derivative-free method

f(R) is not differentiable everywhere...

- ... a derivative-free method
- ... a global search technique

f(R) has many local minima...

- ... a derivative-free method
- ... a global search technique
- ... a fast convergence rate

That was the point!

Our idea: using an hybrid method

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- Use a global exploration component: genetic algorithm (GA)
- Speed up convergence using a local exploitation algorithm Nelder-Mead simplex algorithm (NM)
 - GA alone would be very slow to converge (GA more suitable for discrete search spaces)
 - NM alone would be stuck in local minima (even with restarts)

Coming up next...

The problem: minimum-volume OBB

Bringing optimization into the game

How to solve an optimization problem?

Genetic algorithms (GA)

The Nelder-Mead algorithm (NM)

HYBBRID: let's mix GA and NM together!

Results

Conclusion

Global exploration: genetic algorithms

Stochastic population-based evolutionary method

(original variant proposed by Holland in the 1970s)

- Population-based: keep a large set of candidates at each iteration
- Evolutionary: generate new candidates by combining current ones depending on their performance

Start with a set of candidates (population) and a performance function (fitness function)

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- Selection: parents are selected depending on their fitness
- Crossover: selected parents produce offsprings
- Mutation: offsprings can be subject to mutations
 (random modification, gradient step, SA step, ...)

The Nelder-Mead simplex algorithm

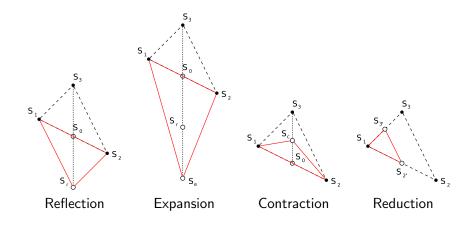
Derivative-free simplicial optimization method

(original algorithm proposed by Nelder & Mead in 1965)

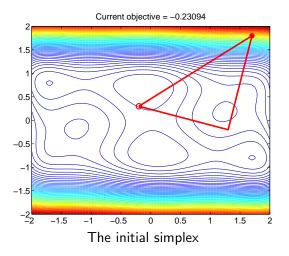
```
Simplex (in \mathbb{R}^n) = set of n+1 affinely independent points
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- $\diamond \ n=2$: triangle
- \diamond n=3: tetrahedron
- ٥ ...

Ideas of the algorithm (details omitted)

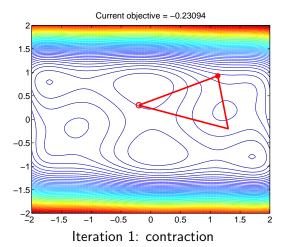


Four main ways to move/transform the simplex depending on the performance of its vertices: affine combinations

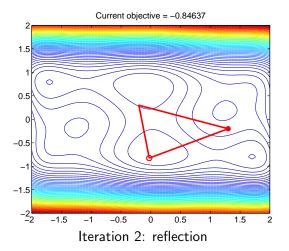




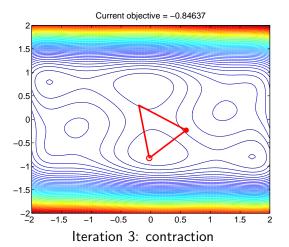


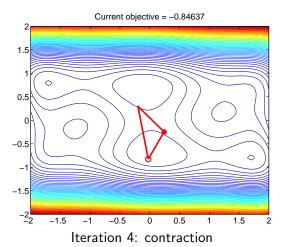


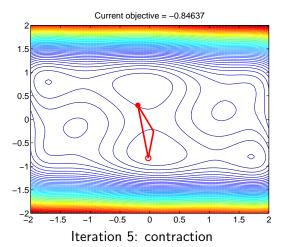




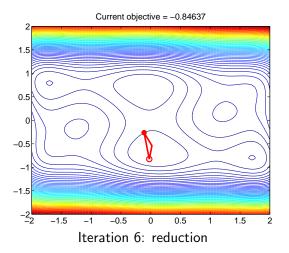




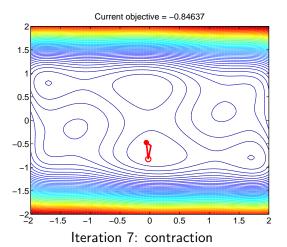




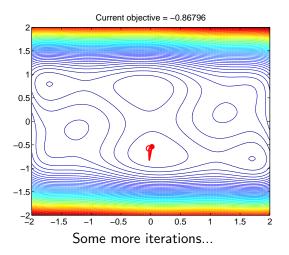




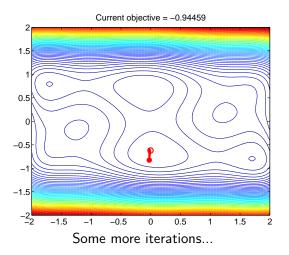




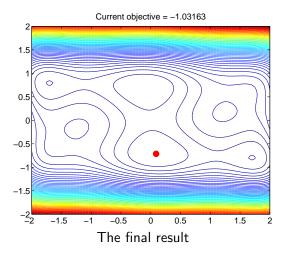
















Population of M simplices (simplex = set of 4 rotation matrices) Fitness function f(R) (volume of corresponding OBB)

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- ♦ Mutation: Apply K Nelder-Mead iterations on each offspring

HYBBRID

Nelder-Mead algorithm \oplus Genetic algorithm on the special orthogonal group SO(3) to solve the optimal OBB problem

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_

HYbrid Bounding Box Rotation IDentification algorithm

Coming up next...

The problem: minimum-volume OBB

Bringing optimization into the game

How to solve an optimization problem?

Results

Behaviour of HYBBRID Comparison to other algorithms

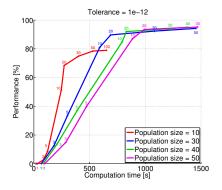
Conclusion

Behavior of HYBBRID

All algorithms tested on a benchmark set of ~ 300 objects (Gamma db) Implementations done in MATLAB (built-in functions are used)

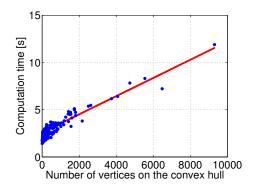
Behavior of HYBBRID: yes, it works!

All algorithms tested on a benchmark set of ~ 300 objects (Gamma db) Implementations done in MATLAB (built-in functions are used)



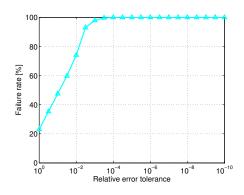
Error is less than 10^{-12} in 90%+ of the cases!

How does it scale?

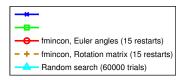


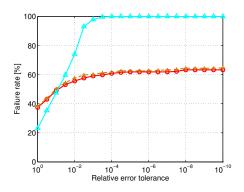
Experimental results show a roughly linear complexity!



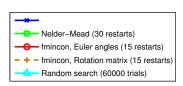


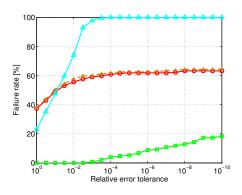
Trying random orientations does not work...



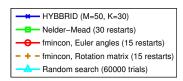


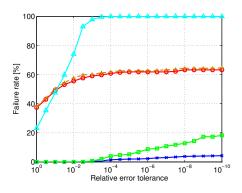
Constrained smooth opti: success rate $\sim 40\%$ but mainly AABBs...





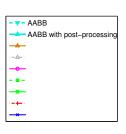
Unconstrained non-diff. opti, random initializations: much better results!

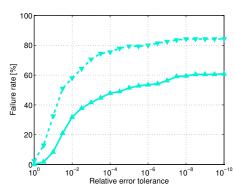




HYBBRID: combining potential solutions does improve the success rate!

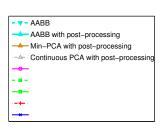
First, let's ignore the computational cost and look at the failure rates...

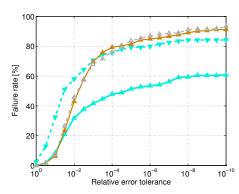




A reference point: the simple AABB

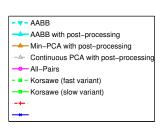
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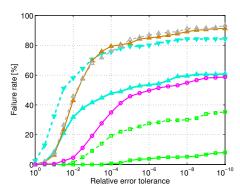




PCA-based methods: limited accuracy

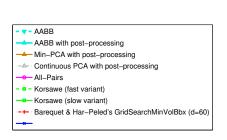
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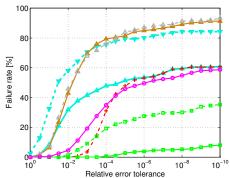




Brute-forcing on a set of orientations may be OK... if well chosen!

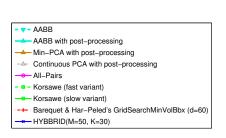
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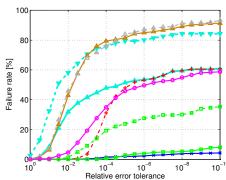




Guaranteed approximation algorithms: limited by computational resources

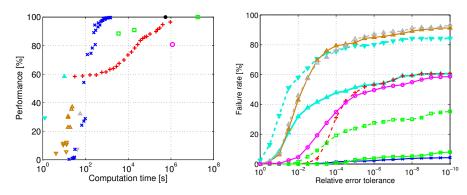
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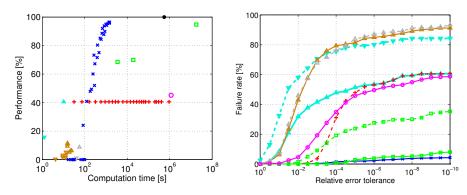
HYBBRID: more accurate than these other methods

What are the computation times? (Tolerance: 10^{-3})



AABB – PCA – Continuous PCA – Brute-force on a set of orientations O'Rourke's exact algorithm – Guaranteed approximation – HYBBRID

What are the computation times? (Tolerance: 10^{-6})



AABB – PCA – Continuous PCA – Brute-force on a set of orientations O'Rourke's exact algorithm – Guaranteed approximation – HYBBRID

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Thank you for your attention!