

A unified conic formulation for convex problems involving powers

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Overview

1. Motivation

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- ◇ Future plans

Convex optimization

Nonlinear optimization

$\min_{x \in \mathbb{R}^n} f_0(x)$ such that $f_i(x) \leq 0$ for all $i \in I$ and $f_i(x) = 0$ for all $i \in E$

- ◇ Variables: finite-dimensional vector $x \in \mathbb{R}^n$
- ◇ Constraints: finite number of (in)equalities, indexed by sets I and E

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Problem is **convex** when

- ◇ objective function f_0 is convex
- ◇ functions f_i defining inequalities $f_i(x) \leq 0$ are convex for all $i \in I$
- ◇ functions f_i defining equalities $f_i(x) = 0$ are affine for all $i \in E$

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 f_0 and f_i are **convex quadratic** for all $i \in I$

$$f_i(x) = x^T Q_i x + r_i^T x + s_i \text{ with } Q_i \succeq 0$$

(equalities f_i , if present, must still be affine for $i \in E$)

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- Convex quadratic can be rewritten using composition of **squared Euclidean norm** and **linear** (vector) function:

$$f_i(x) = \|A_i x\|^2 + (r_i^T x + s_i) \text{ with } Q_i = A_i^T A_i$$

More classes of convex problems

- ◇ Geometric optimization (GO):

f_0 and f_i are **posynomials** (in exponential form) for all $i \in I$

$$f_i(x) = c_i + \sum_{j \in M_i} \exp(a_{ij}x - b_{ij})$$

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- ◇ l_p -norm optimization (l_p O):

f_0 linear, f_i are affine plus sum of convex **powers** with **affine scalar** arguments for all $i \in I$

$$f_i(x) = a_{i0}x - b_{i0} + \sum_{j \in M_i} |a_{ij}x - b_{ij}|^{p_{ij}} \text{ with } p_{ij} \geq 1$$

◇ Sum-of-norm optimization (SNO):

f_0 (and f_i for all $i \in I$, if any) are **convex norms** with affine arguments

$$f_i(x) = \sum_{j \in M_i} \|A_{ij}x - b_{ij}\|_{p_{ij}} \quad \text{with } p_{ij} \geq 1$$

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- ◇ Analytic centering (AC):

f_0 is a sum of **logarithmic** terms, f_i are affine for all $i \in I \cup E$

$$f_0(x) = - \sum_{j \in M_j} \log(a_{ij}x - b_{ij})$$

A common formulation

All above-mentioned problems can be described as follows:

$$\min_{x \in \mathbb{R}^n} f_0(x) \text{ such that } f_i(x) \leq 0 \text{ for all } i \in I \text{ and } f_i(x) = 0 \text{ for all } i \in E$$

with functional terms f_0 and f_i defined by

$$f_i(x) = \sum_{j \in M_i} g_{ij}(A_{ij}x - b_{ij})$$

where **nonlinearity** is confined to functions g_{ij} :

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- ◇ $x \mapsto x$ (identity): for linear optimization and for all **equalities**
- ◇ $x \mapsto \|x\|^2$: for quadratically constrained quadratic optimization
- ◇ $x \mapsto e^x$ for geometric optimization
- ◇ $x \mapsto |x|^p$ with $p \geq 1$ for l_p -norm optimization
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In summary: **separable** functions with every term being the composition of a **simple** (often scalar) **convex nonlinear** function with an **affine** function

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Any algorithm or solver applied to a convex problem will **automatically** benefit from those features

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(in most of the cases: an interior-point method based on the theory of self-concordant barriers)
- ◇ To use those, **additional work** is needed for **each** problem class!
- ◇ Need to exploit **specific structure** of each problem class
- ◇ **Reward** for additional work is better understanding and ability to solve problems more efficiently (including large-scale)

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- ◇ Set K has to be a **convex cone** for useful properties of ordering to hold (and also: closed, solid and pointed for technical reasons)
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- ◇ Conic optimization problems are clearly **convex**
- ◇ **Any convex problem** can be cast as a conic optimization problem
- ◇ The point of a conic formulation is to make it easier to benefit from **active** features of convex optimization (duality and algorithms)

Combining several cones

Considering **several conic** constraints

$$A_1^T y \preceq_{K_1} c_1 \text{ and } A_2^T y \preceq_{K_2} c_2$$

which are equivalent to

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one introduces the **Cartesian product** cone $K = K_1 \times K_2$ to write

$$(c_1 - A_1^T y, c_2 - A_2^T y) \in K_1 \times K_2$$

$$\begin{pmatrix} c_1 \\ c_2 \end{pmatrix} - \begin{pmatrix} A_1^T \\ A_2^T \end{pmatrix} y \succeq_{K_1 \times K_2} 0 \Leftrightarrow A^T y \preceq_K c$$

→ for theory, a single cone can be considered without loss of generality

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- ◇ A good barrier is for example a **self-concordant barrier**, i.e.
 $F : \text{int } K \mapsto \mathbb{R}$ satisfying
 - ▶ F is convex and three times differentiable
 - ▶ $F(x) \rightarrow +\infty$ when $x \rightarrow \partial K$
 - ▶ the following **two** conditions hold

$$\begin{aligned}\nabla^3 F(x)[h, h, h] &\leq 2 \left(\nabla^2 F(x)[h, h] \right)^{\frac{3}{2}} \\ \nabla F(x)^T (\nabla^2 F(x))^{-1} \nabla F(x) &\leq \nu\end{aligned}$$

for all $x \in \text{int } C$ and $h \in \mathbb{R}^n$

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for all $x \in \text{int } C$ and $h \in \mathbb{R}^n$

- ◇ Once a good barrier is known, design of a **polynomial-time algorithm** can be completely straightforward (e.g. using standard short or long step path-following algorithm)

Duality for conic optimization

Problem

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admits a nice **symmetrical** dual

$$\min c^T x \text{ such that } Ax = b \text{ and } x \succeq_{K^*} 0$$

based on the notion of **dual cone**

$$K^* = \{z \in \mathbb{R}^n \text{ such that } x^T z \geq 0 \forall x \in K\}$$

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- ◇ **Weak** duality always holds, **strong** duality holds with regularity assumption (existence of a strictly interior point)
- ◇ Only effort involved in determining a dual problem is computing the **dual cone**
- ◇ Potentially allows design of (symmetrical) **primal-dual** algorithms

Examples of conic optimization

Following 3 cones are (by far) most commonly used

1. $K = \mathbb{R}_+$ is the standard ordering, leading to **linear optimization**
2. $K = \mathbb{L}^n$ leads to **second-order cone optimization** (including QCQO)

$$\mathbb{L}^n = \{(x_0, \dots, x_n) \in \mathbb{R}^{n+1} \mid \sqrt{x_1^2 + \dots + x_n^2} \leq x_0\}$$

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3. $K = \mathbb{S}_+^n$ (positive semidefinite matrices) for **semidefinite optimization**
 - ◇ Those cones share additional theoretical properties (symmetric, i.e. **homogeneous** and **self-dual**)
 - ◇ A fourth cone ($K = \{0\}, K^* = \mathbb{R}$) used for **modelling** convenience
 - ◇ **Many problems** from various domains (e.g. mechanical and electrical engineering, finance) can be modelled using these cones

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 - ◇ **Many problems** from various domains (e.g. mechanical and electrical engineering, finance) can be modelled using these cones
 - ◇ **Many solvers** available for problems involving these cones
 - ◇ However, **no solver** seems available (yet) **for any other cone!**

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- ◇ l_p -norm with rational p possible via construction involving **several** \mathbb{L}^n (size of model increases with "complexity" of p)
- ◇ Geometric optimization can only be **approximated** using several \mathbb{L}^n (size of model increases with accuracy required)

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Our aim: a **single** family of convex **cone** to model **all** of these **exactly**

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The power cone \mathcal{P}_p

Let us consider the **epigraph** of the convex power function

$$z \mapsto |z|^p \text{ with } p \geq 1 \quad \rightarrow \quad E_p = \{(x, z) \mid |z|^p \leq x\}$$

and take its **conic hull**: $(x, y, z) \in K_{E_p} \Leftrightarrow \frac{1}{y}(x, z) \in \mathcal{E}_p$

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The resulting 3-dimensional cone will be called the **power cone** and denoted

$$\mathcal{P}_p = \{(x, y, z) \in \mathbb{R}_+ \times \mathbb{R}_+ \times \mathbb{R} \mid x^{\frac{1}{p}} y^{\frac{1}{q}} \geq |z|\}$$

(with the usual convention $\frac{1}{p} + \frac{1}{q} = 1$)

The power cone \mathcal{P}_p

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(with the usual convention $\frac{1}{p} + \frac{1}{q} = 1$)

- ◇ \mathcal{P}_p is equivalent to (rotated) \mathbb{L}^3 when $p = q = 2$
- ◇ \mathcal{P}_p can be obtained from several \mathbb{L}^n for rational values of p (number of cones required increases with "complexity" of p)

Duality

The **dual** of the **power cone**

$$\mathcal{P}_p = \{(x, y, z) \in \mathbb{R}_+ \times \mathbb{R}_+ \times \mathbb{R} \mid x^{\frac{1}{p}} y^{\frac{1}{q}} \geq |z|\}$$

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is ... the **power cone** itself !

(up to a simple scaling of the variables/use of a different inner product)

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Power cone is therefore **self-dual** (but not homogeneous), and actual dual cone (for standard inner product) is

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Therefore the **dual** for any conic problem **based on** \mathcal{P}_p

- ◇ is a conic problem **also based on** \mathcal{P}_p
- ◇ can be derived in a completely **mechanical** way

Barrier function for \mathcal{P}_p

- ◇ A self-concordant barrier for

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with parameter $\nu = 4$ was proposed by NESTEROV

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- ◇ This implies complexity of solving conic problems involving \mathcal{P}_p depends only the number of cones (not on parameter p)

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- ◇ Why convex optimization?
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2. Unified conic formulation

- ◇ The power cone
- ◇ Modelling problems involving powers using the power cone
- ◇ Modelling problems involving exponentials using the power cone

3. Concluding remarks

- ◇ Future plans

Modelling with the power cone

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- For example, modelling the constraint

$$|u_1 + u_2|^3 + |u_1 - u_2|^{4.5} \leq 2u_2 + 1$$

will be done with

$$\begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix} - \begin{pmatrix} 0 & 0 & -1 \\ 0 & 0 & 0 \\ -1 & -1 & 0 \end{pmatrix} \begin{pmatrix} u_1 \\ u_2 \\ v \end{pmatrix} \in \mathcal{P}_3, \quad \begin{pmatrix} 1 \\ 1 \\ 0 \end{pmatrix} - \begin{pmatrix} 0 & -2 & 1 \\ 0 & 0 & 0 \\ -1 & 1 & 0 \end{pmatrix} \begin{pmatrix} u_1 \\ u_2 \\ v \end{pmatrix} \in \mathcal{P}_{4.5}$$

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$$\Leftrightarrow (v, 1, u_1 + u_2) \in \mathcal{P}_3 \text{ and } (1 + 2u_2 - v, 1, u_1 - u_2) \in \mathcal{P}_{4.5}$$

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- Those are examples of constraints where nonlinearity is additively **separable into scalar** convex components

Modelling non-separable nonlinearity

- ◇ Norm constraint

$$\|z\|_p \leq t \quad \Leftrightarrow \quad |z_1|^p + |z_2|^p \cdots |z_n|^p \leq t^p \text{ and } t \geq 0$$

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- ◇ We use the following trick:

$$\begin{aligned} & |z_1|^p + |z_2|^p \cdots |z_n|^p \leq t^p \\ \Leftrightarrow & \left| \frac{z_1}{t} \right|^p + \left| \frac{z_2}{t} \right|^p \cdots \left| \frac{z_n}{t} \right|^p \leq 1 \\ \Leftrightarrow & \left| \frac{z_1}{t} \right|^p \leq \frac{x_1}{t}, \dots, \left| \frac{z_n}{t} \right|^p \leq \frac{x_n}{t} \text{ and } \frac{x_1}{t} + \cdots + \frac{x_n}{t} = 1 \\ \Leftrightarrow & (x_1, t, y_1) \in \mathcal{P}_p, \dots, (x_n, t, y_n) \in \mathcal{P}_p \text{ and } x_1 + \cdots + x_n = t \end{aligned}$$

which can be modelled using conic optimization

Norm constraints (continued)

- ◇ Using the previous construction, we can model any constraint $\|x\|_p \leq 1$ for $p \geq 1$
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- ◇ When $p = 2$, this is the **second-order cone**: we can therefore also model all second-order cone optimization problems, including **QCQO** problems
- ◇ We can also model more complicated non-separable expressions, such as

$$z_1^{2.5} z_2^{-4.5} z_3^4 + 2z_1 z_2 + z_3^4 \leq t^2$$

(crucial condition for convexity is that degree of every term on l.h.s. should be greater than degree of r.h.s.)

- ◇ This looks similar to posynomials involved in geometric optimization but is completely different (no exponential transformation)

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$$\begin{aligned} x^{\frac{1}{p}} y^{\frac{1}{q}} \geq \left| y + \frac{z'}{p} \right| &\Leftrightarrow x^{\frac{1}{p}} y^{\left(\frac{1}{q}-1\right)} \geq \left| 1 + \frac{z'/y}{p} \right| \Leftrightarrow xy^{-1} \geq \left| 1 + \frac{1}{p} \frac{z'}{y} \right|^p \\ &\Leftrightarrow \left| 1 + \frac{1}{p} \left(\frac{z'}{y} \right) \right|^p \leq \frac{x}{y} \end{aligned}$$

Letting now p tend to $+\infty$ we obtain

$$\left|1 + \frac{1}{p} \left(\frac{z'}{y}\right)\right|^p \leq \frac{x}{y} \quad \longrightarrow_{p \rightarrow \infty} \quad \exp\left(\frac{z'}{y}\right) \leq \frac{x}{y}$$

which defines the **exponential cone**:

$$\mathcal{E}_p = \left\{ (x, y, z) \in \mathbb{R}_+ \times \mathbb{R}_+ \times \mathbb{R} \mid \exp\left(\frac{z}{y}\right) \leq \frac{x}{y} \right\}$$

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 and therefore **analytic centering**
- ◇ We can also model the **epigraph of entropy**: $(1, y, -z) \in \mathcal{E}_p \Leftrightarrow \exp(-z/y) \leq 1/y \Leftrightarrow -z/y \leq -\log y \Leftrightarrow y \log y \leq z$ and therefore **entropy optimization**

Barrier for exponential cone

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Is something wrong? Can be corrected by adding a missing **constant term**:

$$\begin{aligned} F_{\text{exp}}(x, y, z) &= \lim_{p \rightarrow +\infty} \left(F_p(x, y, y + z/p) - \log \frac{p}{2} \right) \\ &= \dots = -\log(y \log(x/y) - z) - \log(x) - \log(y) \end{aligned}$$

Using this unified barrier, one can solve any conic problem involving exponential cones \mathcal{E}_p
 (and recompute the standard barriers for exponential, minus logarithm and entropy epigraphs)

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An example: the **Lambert W function**, defined by $W(x) \exp W(x) = x$

From *MathWorld*: Banwell and Jayakumar (2000) showed that a W -function describes the relation between voltage, current and resistance in a diode, and Packel and Yuen (2004) applied the W -function to a ballistic projectile in the presence of air resistance. Other applications have been discovered in statistical mechanics, quantum chemistry, combinatorics, enzyme kinetics, the physiology of vision, the engineering of thin films, hydrology, and the analysis of algorithms (Hayes 2005).

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Indeed, we can check that

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In summary, combining a quadratic and an exponential constraint, we have shown that

$$0 \leq y \leq W(x) \Leftrightarrow (x, y, z) \in \mathcal{E}_p \text{ and } (z, 1, x) \in \mathcal{P}_2$$

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Potential drawback: conic modelling sometimes require the introduction of a large number of **additional variables** (e.g. $\|x\|_p \leq t$ constraint)

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joint work with Robert Chares, CORE

see talk later today, session TC4 at 2:45pm

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- ◇ In the future: take advantage of self-duality and implement a primal-dual interior-point method

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- ◇ **Stay tuned !**
- ◇ If you have any problem that you think could be modelled in our framework, please do not hesitate to **contact us!**

Thank you for your attention!