FORWARD/BACKWARD DECOMPOSITION OF PERIODIC DESCRIPTOR SYSTEMS

J. Sreedhar*

Department of Electrical & Computer Engineering University of Illinois at Urbana-Champaign, IL 61801-2307, USA. Tel number: 1-217-333-7541

Paul Van Dooren[†]

Department of Mathematical Engineering, Université Catholique de Louvain, 1348 Louvain-la-Neuve, Belgium. Tel number: 32-10-47-8040

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Abstract

We propose a Forward/Backward (F/B) decomposition of periodic discrete time descriptor systems and describe a numerically reliable method to compute it. We mention a couple of applications of this F/B canonical form; namely, the solution of two-point boundary value problems, and discrete time Floquet theory.

1 Introduction

We consider the periodically time varying (PTV) system

$$\Sigma: \quad E(k)x(k+1) = A(k)x(k), \quad k \in \mathbb{Z}^+, \tag{1}$$

where \mathbb{Z}^+ is the set of non-negative integers, x(k) is an n-dimensional vector of descriptor variables, and K is the smallest positive integer for which E(k) = E(k + K), A(k) = A(k + K), $\forall k \in \mathbb{Z}^+$. We assume that the matrices E(k) and A(k) are real, $n \times n$, and satisfy the

condition

$$\lambda \begin{bmatrix} E_0 & & & & & \\ & E_1 & & & \\ & & \ddots & & \\ & & & E_{K-1} \end{bmatrix} - \begin{bmatrix} A_1 & & & & A_0 \\ & \ddots & & & \\ & & A_{K-1} & & \end{bmatrix}$$
(2)

is regular.

Recently [1] we have shown that statement (2) is equivalent to *solvability* or *conditionability* of Σ , an assumption commonly made in connection with the Two-point Boundary Value Problem (TPBVP)

Here N simply specifies a time interval of interest. For a discussion of the general time-varying TPBVP, see Luenberger [2, 3]. When studying equation (1), it is common to write

$$V(k) E(k) W(k+1) x(k+1) = V(k) A(k) W(k) x(k),$$
(4)

where V(k) = V(k + K), W(k) = W(k + K) are nonsingular for all k, and try to achieve as much simplification

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in form as possible for E(k), A(k). For instance, in the time-invariant (TI) case, where none of the above matrices depend on k, one could reduce E, A to Schur form [4] or decompose the system into "forward" and "backward" subsystems [5]. In this note, for the periodic case, we take the latter approach and demonstrate that it is possible to obtain a canonical forward/backward (F/B) form in (4):

$$E(k) \leftarrow \begin{bmatrix} I & 0 \\ 0 & E_b(k) \end{bmatrix}, \quad A(k) \leftarrow \begin{bmatrix} A_f(k) & 0 \\ 0 & I \end{bmatrix},$$

where the subscripts 'f' and 'b' refer to forward and backward components. Such a F/B form appears to be quite useful in the general study of periodic systems, e.g., in realization [6]. As an immediate application, we demonstrate that the decomposition into F/B form greatly simplifies the boundary recursion formulae of Luenberger [3] and also helps to extend an earlier result in discrete time Floquet theory [7] to descriptor systems.

2 A spectral decomposition theorem and the F/B form

We begin by introducing some notation. For Σ , it is well-known that the spectrum, i.e., the set of eigenvalues counting multiplicities, of the periodic pencil $\lambda E(k) - A(k)$ plays an important role; we shall denote it by $\lambda([E, A])$. The special situations where either $E(k) = I_n$ or $A(k) = I_n$ afford some simplification and we mention these separately. When $E(k) = I_n$, $k \in \mathbb{Z}^+$, we have the forward transition matrix

$$\phi_{A}(k, k_{0}) = \begin{cases}
A(k-1)A(k-2) \dots A(k_{0}+1)A(k_{0}) & \text{if } k > k_{0} \geq 0, \\
I & \text{if } k = k_{0} \geq 0, \\
\phi_{A}(k, k_{0}) & \text{undefined for } k < k_{0},
\end{cases}$$
(5)

and the characteristic multipliers $\lambda_f([A]) := \lambda([I, A])$, which are nothing but the eigenvalues of the monodromy matrix $\phi_A(k_0 + K, k_0)$. Similarly, when $A(k) = I_n$, $k \in \mathbb{Z}^+$, we have the backward transition matrix

$$\varphi_{E}(k_{0}, k) = \begin{cases}
E(k_{0})E(k_{0} + 1) \dots E(k - 1)E(k - 1) & \text{if } k > k_{0} \geq 0, \\
I & \text{if } k = k_{0} \geq 0, \\
\varphi_{E}(k_{0}, k) & \text{undefined for } k < k_{0},
\end{cases}$$
(6)

and the characteristic multipliers $\lambda_b([E]) := \lambda([E, I])$, which are nothing but the eigenvalues of the monodromy matrix $\varphi_E(k_0, k_0 + K)$.

We now state a spectral decomposition theorem, which extends a well known result for time invariant pencils to the periodic case.

Theorem 1. Given E(k), A(k) satisfying (1), (2), and an arbitrary disjoint partition Λ_1 , Λ_2 of the spectrum $\lambda([E, A])$, i.e., $\Lambda_1, \Lambda_2 \subseteq \lambda([E, A])$, $\Lambda_1 \cup \Lambda_2 = \lambda([E, A])$, $\Lambda_1 \cap \Lambda_2 = \emptyset$, it is always possible to find K-periodic nonsingular matrices V(k) and W(k) such that

$$V(k) E(k) W(k+1) = \begin{bmatrix} E_{11}(k) & 0 \\ 0 & E_{22}(k) \end{bmatrix},$$
$$V(k) A(k) W(k) = \begin{bmatrix} A_{11}(k) & 0 \\ 0 & A_{22}(k) \end{bmatrix},$$

where all matrices on the right are upper-triangular, and

$$\lambda([E_{11}, A_{11}]) = \Lambda_1, \quad \lambda([E_{22}, A_{22}]) = \Lambda_2.$$

Sketch of proof. Step 1. Put E(k) and A(k) in upper-triangular form using the periodic Schur algorithm [8], which always exists (and defines the spectrum unambiguously when (2) is satisfied). Arrange, by reordering diagonal elements if necessary, that

$$E(k) = \begin{bmatrix} E_{11}(k) & E_{12}(k) \\ 0 & E_{22}(k) \end{bmatrix}, \quad A(k) = \begin{bmatrix} A_{11}(k) & A_{12}(k) \\ 0 & A_{22}(k) \end{bmatrix},$$

with $\lambda([E_{11}, A_{11}]) = \Lambda_1$ and $\lambda([E_{22}, A_{22}]) = \Lambda_2$. This step uses unitary periodic transformations only.

Step 2. Since $\lambda([E_{11}, A_{11}]) \cap \lambda([E_{22}, A_{22}]) = \emptyset$, zero each off-diagonal block using a further (nonunitary) periodic transformation. This involves solving the periodic Sylvester equation

$$\begin{bmatrix} I & L(k) \\ 0 & I \end{bmatrix}^{-1} \begin{bmatrix} E_{11}(k) & E_{12}(k) \\ 0 & E_{22}(k) \end{bmatrix} \begin{bmatrix} I & R(k+1) \\ 0 & I \end{bmatrix}$$

$$= \begin{bmatrix} E_{11}(k) & 0 \\ 0 & E_{22}(k) \end{bmatrix}$$

$$\begin{bmatrix} I & L(k) \\ 0 & I \end{bmatrix}^{-1} \begin{bmatrix} A_{11}(k) & A_{12}(k) \\ 0 & A_{22}(k) \end{bmatrix} \begin{bmatrix} I & R(k) \\ 0 & I \end{bmatrix}$$

$$= \begin{bmatrix} A_{11}(k) & 0 \\ 0 & A_{22}(k) \end{bmatrix}.$$

The left and right transformations from Steps 1 and 2 can be combined to get V(k) and W(k).

Remark 1. By choosing Λ_1, Λ_2 appropriately in Theorem 1, it is possible to have $\lambda([E_{11},A_{11}])$ lie on or inside the unit circle and $\lambda([E_{22},A_{22}])$ outside the unit circle. Note that $E_{11}(k)$ and $A_{22}(k)$ are then nonsingular. Such a form would be useful, for instance, in finding a minimal order TI description of Σ using strict system equivalence [9] operations on its stacked [10] representation. This work is currently under progress [11]. See [12] for details of this procedure for standard state space systems, viz., when $E(k) = I_n, k \in \mathbb{Z}^+$.

Remark 2. Theorem 1 would be useful when one considers the additive decomposition problem for periodic discrete-time systems. See [13] for details in the time-invariant case

Remark 3. Theorem 1 shows the possibility of diagonalizing E(k) and A(k) when the eigenvalues $\lambda([E, A])$ are distinct.

Theorem 1 quickly leads to the promised F/B form, as shown next.

Corollary 1. Given E(k) and A(k) satisfying (1), (2), it is always possible to find K-periodic nonsingular matrices V(k) and W(k) such that

$$V(k) E(k) W(k+1) = \begin{bmatrix} I & 0 \\ 0 & E_b(k) \end{bmatrix},$$

$$V(k) A(k) W(k) = \begin{bmatrix} A_f(k) & 0 \\ 0 & I \end{bmatrix},$$
(7)

where $E_b(k)$ and $A_f(k)$ are upper-triangular, $\lambda_f([A_f])$ lies on or inside the unit circle, and $\lambda_b([E_b])$ lies inside the unit circle.

Sketch of proof. Use Theorem 1 (in particular, see Remark 1). Premultiply with $\operatorname{diag}(E_{11}^{-1}(k), A_{22}^{-1}(k))$ to get the final result in (7), with

$$A_f(k) = E_{11}^{-1}(k)A_{11}(k), \quad E_b(k) = A_{22}^{-1}(k)E_{22}(k). \quad \Box$$

We conclude this note by mentioning two uses of the F/B canonical form, namely, periodic TPBVPs and discrete time Floquet theory.

3 Two-point boundary value problems

Clearly, Corollary 1 decomposes the original system Σ into a forward and a backward part:

$$x_f(k+1) = A_f(k)x_f(k), \quad A_f(k+K) = A_f(k),$$
(8f)

$$E_b(k)x_b(k+1) = x_b(k), \quad E_b(k+K) = E_b(k),$$
 (8b)

Given $x_f(0)$, we can iterate (??f) forward in time to obtain $x_f(k)$ for all $k \geq 0$, and, similarly, given $x_b(N)$, we can iterate (??b) backwards in time to obtain $x_b(k)$ for all $0 \leq k \leq N$. The solution will, of course, involve the forward and backward transition matrices introduced in (5) and (6).

Alternatively, if the boundary conditions are specified as in (3), then the F/B decomposition yields very simple expressions for Luenberger's boundary recursion procedure [3] applied to Σ , as we show next. The idea behind the procedure is to express the inherent linear restrictions on x(0) and x(N) in the form of a boundary mapping equation $Z_0(0, N)x(0) + Z_N(0, N)x(N) = 0$. Likewise, one could write

$$Z_0(0,k)x(0) + Z_k(0,k)x(k) = 0, \quad k = 1, 2, \dots, N,$$
 (9)

where $Z_0(0, k)$ and $Z_k(0, k)$ are some appropriate $n \times n$ matrices, to make explicit the linear restrictions on x(0) and x(k) that are imposed implicitly by (1). The basic step is boundary recursion, wherein $Z_0(0, k)$ and $Z_k(0, k)$ are found recursively as follows:

$$\begin{bmatrix}
F(k) & G(k) \\
H(k) & J(k)
\end{bmatrix}
\begin{bmatrix}
Z_0(0, k) & Z_k(0, k) & 0 \\
0 & -A(k) & E(k)
\end{bmatrix}$$

$$= \begin{bmatrix}
Z_0(0, k+1) & 0 & Z_{k+1}(0, k+1) \\
X(k) & I & Y(k)
\end{bmatrix}. (10)$$

It is easy to see that F(k) = A(k), $G(k) = Z_k(0, k)$ will produce the desired zero block on the right hand side of (10). Furthermore, when E(k) and A(k) are in the F/B form, we have

Comparing (10) with (11), we get all the necessary formulae for boundary recursion of periodic systems. Finally, the boundary conditions W(0)x(0) + W(N)x(N) = w, with $W(0) = [W_{0f}, W_{0b}], W(N) = [W_{Nf}, W_{Nb}]$ partitioned according to (8), are independent of the boundary mapping of Σ if

$$[W_{0f} + \phi_{A_f}(N, 0)W_{Nf}, W_{Nb} + \varphi_{E_b}(0, N)W_{0b}]$$
 (12)

is nonsingular. Equations (11) and (12) are analogous to the TI case discussed in [3].

4 Discrete time Floquet theory

Since time invariant linear system theory is much more developed than its periodic counterpart, it is natural to wonder under what conditions can the PTV system Σ be converted into a TI system by a nonsingular transformation as shown in equation (4). A transformation of the kind

$$E(k) \leftarrow Z(k) E(k) Q(k+1), \quad A(k) \leftarrow Z(k) A(k) Q(k),$$

where the matrices Q(k), Z(k) are nonsingular for all k, is known as a *kinematic similarity*, and a K-periodic similarity when Q(k), Z(k) are, in addition, periodic of period K.

In this section, we give a necessary and sufficient condition for a K-periodic pair $\{A(k), E(k)\}$ to be K-periodically similar to a constant pair $\{E, A\}$.

We need a couple of lemmas first. Lemma 1 states when a forward PTV system

$$x(k+1) = A_f(k)x(k), \quad A_f(k+K) = A_f(k),$$

can be transformed into a TI system

$$\hat{x}(k+1) = A_f \hat{x}(k)$$

under an invertible, K-periodic, change of variable $\hat{x}(k) := T_f^{-1}(k)x(k)$. See also [7], [14].

Lemma 1. Given K-periodic matrices $A_f(k)$, there exist K-periodic invertible matrices $T_f(k)$ satisfying

$$T_f^{-1}(k+1)A_f(k)T_f(k) = A_f, \quad k \in \mathbb{Z}^+,$$
 (13)

if and only if

nullity
$$\begin{bmatrix} A_f(j+i-1) & & & & \\ I & \ddots & & & \\ & \ddots & & & \\ & & I & A_f(j+1) & \\ & & I & A_f(j) \end{bmatrix}$$

is independent of j, for $1 \le i \le n$, $1 \le j \le K$. (14)

Proof. In [7], we proved that (13) is satisfied if and only if

nullity
$$\{A_f(j+i-1)\cdots A_f(j+1)A_f(j)\}$$

is independent of j, for 1 < i < n, 1 < j < K. (15)

But $A_f(j+i-1)\cdots A_f(j+1)A_f(j)=\phi_A(j+i,j)$ has the same nullity as the matrix in (14), because the latter matrix is, modulo right and left elementary operations, nothing but

diag
$$\{I,\ldots,I,\phi_A(j+i,j)\}.$$

Therefore, (14) is the same as (4).

We state the next lemma without proof; it gives a corresponding result for the backward PTV system

$$E_b(k)x(k+1) = x(k), \quad E_b(k+K) = E_b(k).$$

Lemma 2. Given K-periodic matrices $E_b(k)$, there exist K-periodic invertible matrices $T_b(k)$ satisfying

$$T_b^{-1}(k)E_b(k)T_b(k+1) = E_b, \quad k \in \mathbb{Z}^+,$$
 (16)

if and only if

nullity
$$\begin{bmatrix} E_b(j+i-1) & I & & & \\ & \ddots & \ddots & & \\ & & E_b(j+1) & I \\ & & & E_b(j) \end{bmatrix}$$

is independent of j, for $1 \le i \le n$, $1 \le j \le K$. (17)

We now state the conditions under which Σ is K-periodically similar to a TI descriptor system. The F/B form helps in deriving this result.

Theorem 2. Given K-periodic matrices E(k), A(k), there exist K-periodic invertible matrices Z(k), Q(k) satisfying

$$Z(k) E(k) Q(k+1) = E, \quad Z(k) A(k) Q(k) = A, \quad k \in \mathbb{Z}^+,$$
(18)

if and only if the nullities of

$$\begin{bmatrix} A(j+i-1) \\ E(j+i-2) & \ddots & & \\ & & \ddots & \\ & & A(j+1) \\ E(j) & A(j) \end{bmatrix}$$
and
$$\begin{bmatrix} E(j+i-1) & A(j+i-1) & & \\ & & \ddots & & \\ & & E(j+1) & A(j+1) \\ E(j) & & E(j) \end{bmatrix}$$
(19)

are independent of j, for $1 \le i \le n$, $1 \le j \le K$.

Sketch of proof. Necessity is easy to prove because (18) implies

$$\begin{bmatrix} \ddots & & & & & \\ & Z(j+i-1) & & & & \\ & & Z(j) \end{bmatrix}$$

$$\begin{bmatrix} \ddots & & \ddots & & & \\ & E(j+i-1) & A(j+i-1) & & & \\ & & \ddots & & & \\ & & E(j) & A(j) \end{bmatrix}$$

$$\begin{bmatrix} \ddots & & & & & \\ & E(j) & A(j) \end{bmatrix}$$

$$\begin{bmatrix} \ddots & & & & \\ & Q(j+i) & & \\ & & \ddots & \\ & & & E & A \end{bmatrix}$$

$$\begin{bmatrix} \ddots & & & & \\ & E & A & \\ & & \ddots & \\ & & E & A \end{bmatrix}$$

$$(20)$$

for all $j \in \mathbb{Z}^+$. To show sufficiency, we must invoke the F/B form and *perfect shuffle*, followed by Lemmas 1 and 2, We omit the details here due to lack of space. \square

5 Conclusion

In this note, we studied spectral decomposition of periodic discrete time descriptor systems and described a

numerically reliable method to compute it. This led to a Forward/Backward (F/B) form which is useful for both numerical computation and theory—as shown for instance in Sections 3 and 4, respectively. We expect that further applications for these decompositions will be found.

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